

∞ -categories w/ adjoints

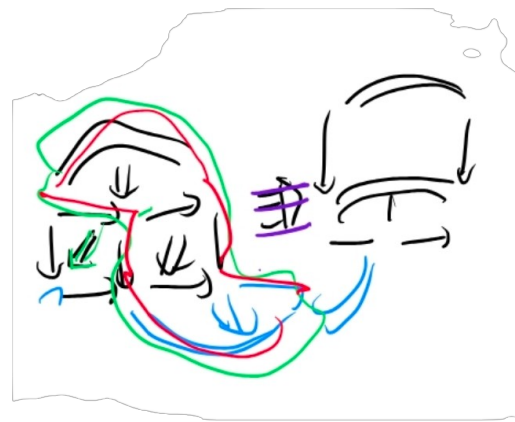
are a bit like spaces

$$\downarrow \otimes \rightarrow = \downarrow$$

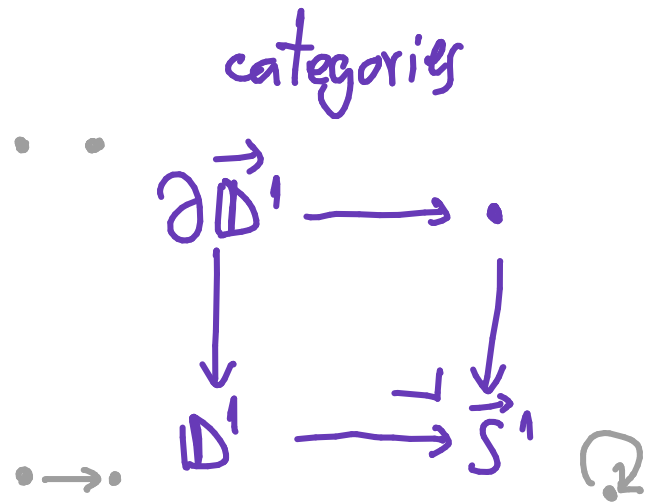
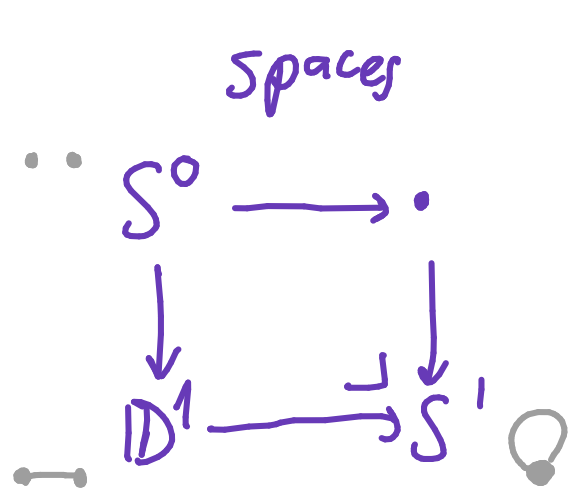
Daniel Teixeira, Dalhousie University

CMS 2026

$$\dots F^{\cup} \dashv F^{\cup L} \dashv F \dashv F^{\cup R} \dashv F^{\cup RR} \dashv \dots$$



Program: directed homotopy theory $\sim (\infty, \infty)$ -categories



$$\pi_1(X) = \text{Map}_{S^1} (S^1, (X, x)) \sim$$

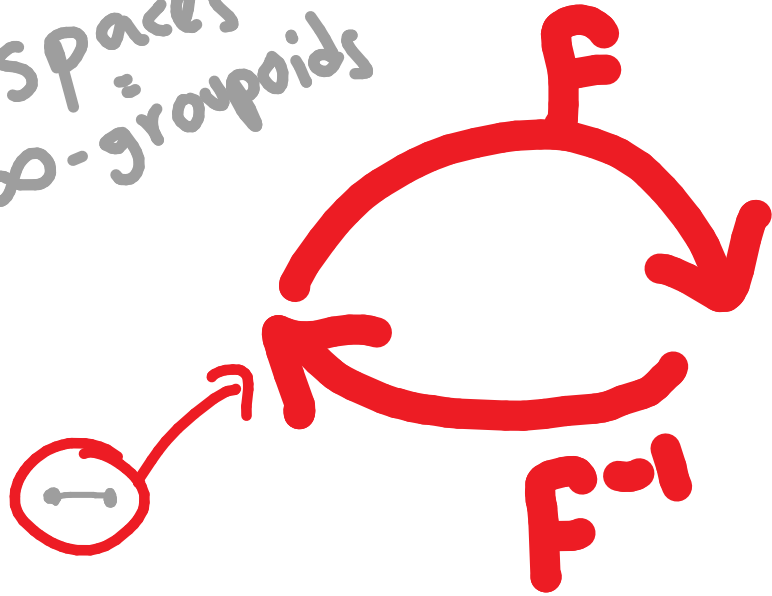
\hookrightarrow loops

$$\pi_1(\vec{S}^1, (C, c)) = \text{Map}_{\text{Cat}_1} (\vec{S}^1, (C, c))$$

\hookrightarrow directed loops

Program: directed homotopy theory $\sim (\infty, \infty)$ -categories

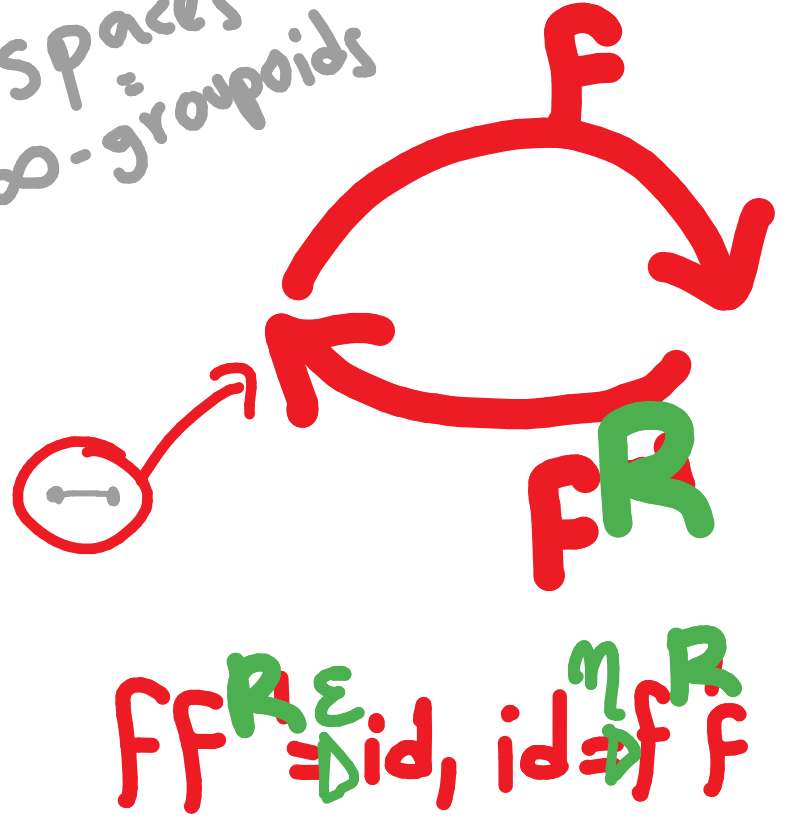
spaces
 ∞ -groupoids



$$FF^{-1} = \text{id}, \text{id} = F^{-1}F$$

Program: directed homotopy theory $\sim (\infty, \infty)$ -categories

spaces
 ∞ -groupoids

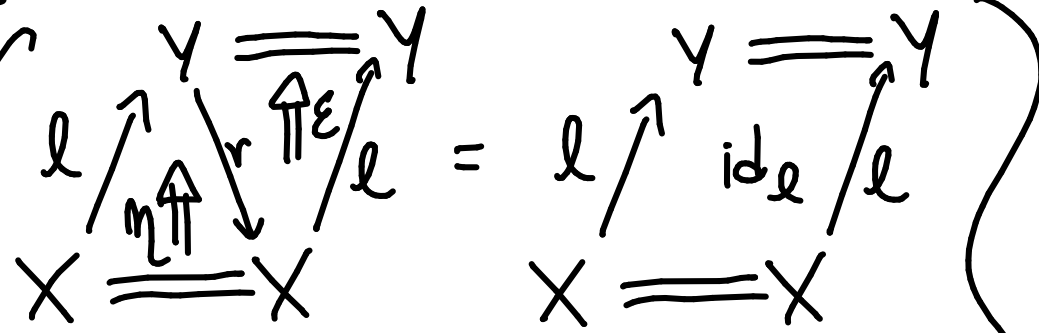


maybe?

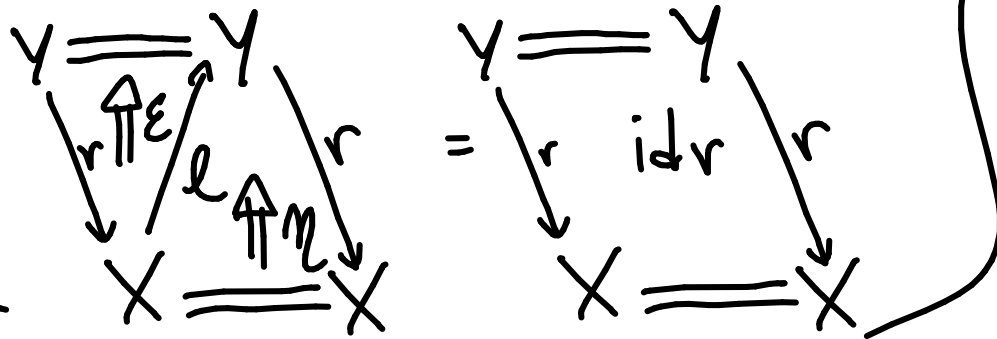
- connections with dualizability
- TQFTs
- weak n -categories

Adjunctions in 2-categories

Def an adjunction in $\mathcal{C} \in 2\text{Cat}$



Def adj = walking adjunction =

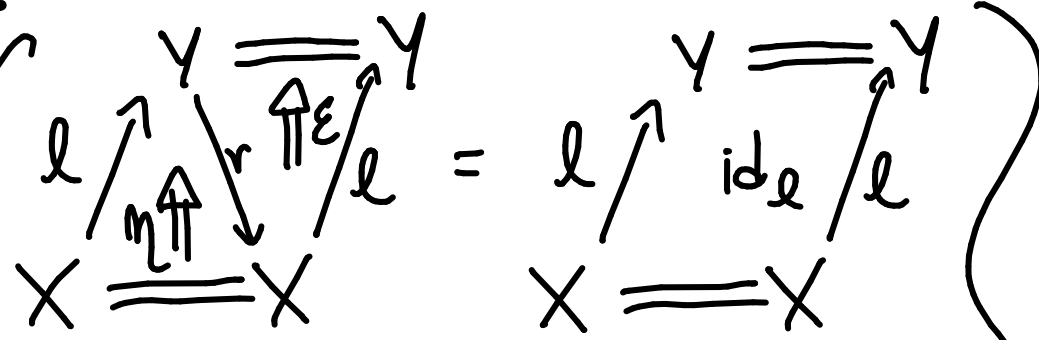


Rmk $\uparrow \xrightarrow{l \text{ or } r} \text{adj}$
is epi...

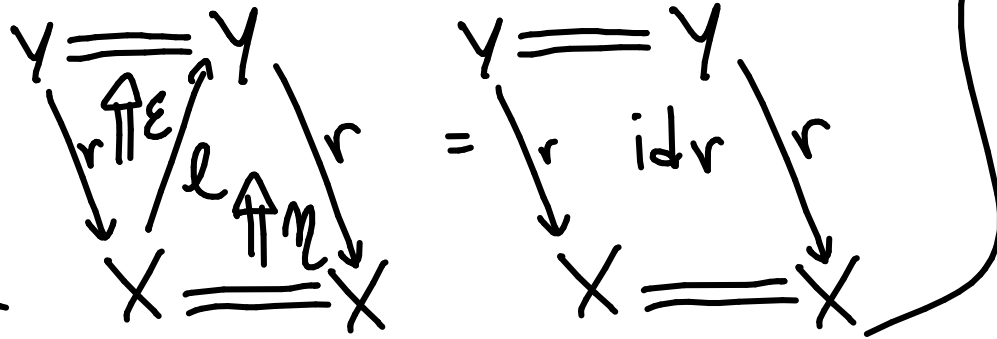
rmk: I never use the word "pseudo"

Adjunctions in 2-categories

Def an adjunction in $\mathcal{C} \in 2\text{Cat}$



Def adj = walking adjunction =



Rmk \rightarrow is epi $\xrightarrow{l \text{ or } r}$ adj

rmk: I never use the word "pseudo"

Adjunctions in $(\infty, 2)$ -categories

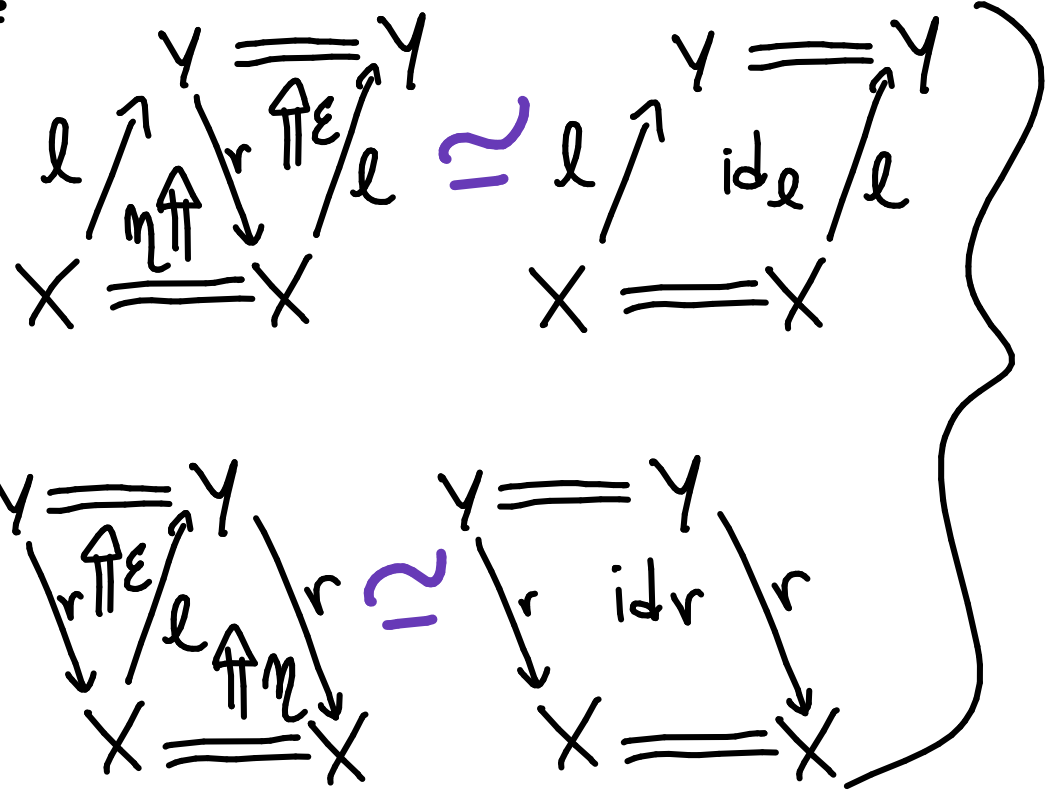
$$(\Rightarrow) \circ (\Leftarrow) \xrightarrow{\sim} (\Leftarrow) \circ (\Rightarrow)$$

Def? an adjunction in $\mathcal{C} \in (\infty, 2)\text{Cat}$

There is such a thing. But

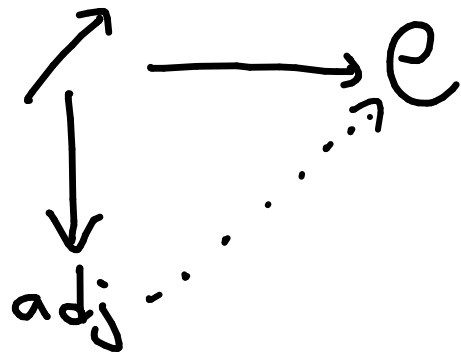
Thm $\text{adj} \xrightarrow{\sim} \text{hadj}$,
(Riehl-Verity, 16')

so we just use $\text{adj} \in 2\text{Cat}$
 \downarrow
 $(\infty, 2)\text{Cat}$



Adjunctions in (∞, ∞) -categories

1-morphisms



suspense

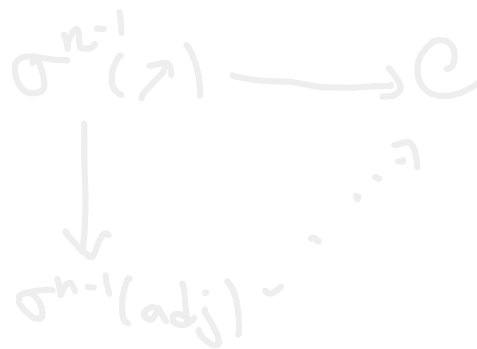
$$\begin{array}{ccc} \infty\text{Cat} & \xrightarrow{\sigma} & \infty\text{Cat} \\ \mathcal{C} & \xrightarrow{\quad} & \mathcal{C} \end{array}$$

(ex) $\mathcal{C} \in \text{Cat} \rightarrow \sigma \mathcal{C} \in \text{Cat}$

(ex) $\vec{D}^n \sim \vec{D}^{n+1} := \sigma \vec{D}^{n+1}$

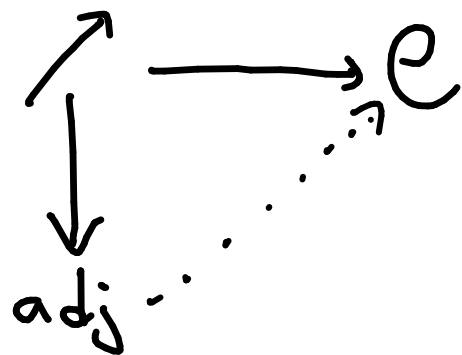
$\downarrow \sim \circlearrowleft \downarrow \circlearrowright$

n -morphisms



Adjunctions in (∞, ∞) -categories

1-morphisms



suspense

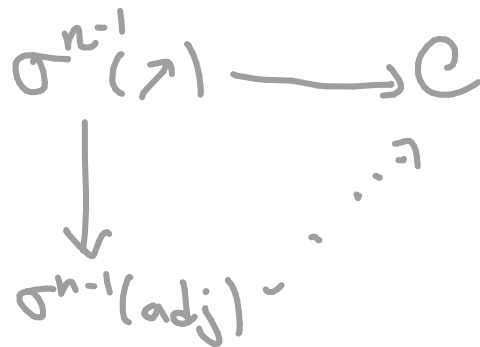
$$\begin{array}{ccc} \infty\text{Cat} & \xrightarrow{\sigma} & \infty\text{Cat} \\ \mathcal{C} & \xrightarrow{\quad} & \mathcal{C} \end{array}$$

(ex) $\mathcal{C} \text{enCat} \rightarrow \sigma \mathcal{C} \text{en}(n+1) \text{Cat}$

(ex) $\vec{D}^n \sim \vec{D}^{n+1} := \sigma D^{n+1}$

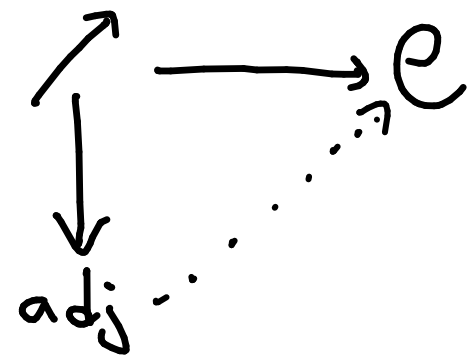
$\downarrow \sim \circlearrowleft \downarrow \circlearrowright$

n -morphisms



Adjunctions in (∞, ∞) -categories

1-morphisms



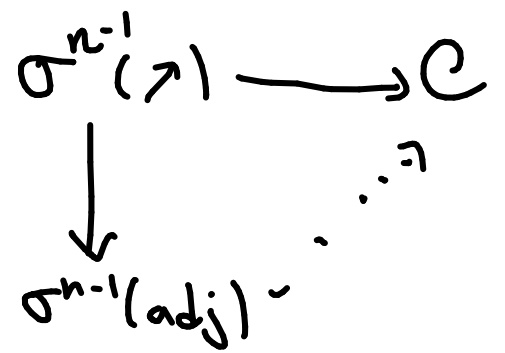
suspense

$$\infty\text{Cat} \xrightarrow{\sigma} \infty\text{Cat}$$

$$e \longmapsto \bullet \xrightarrow{e} \bullet$$

- ex $\text{CenCat} \rightarrow \sigma \text{CenCat}$
 - ex $\vec{D}^n \sim \vec{D}^{n+1} := \sigma \vec{D}^{n+1}$
- $$\downarrow \sim \circlearrowleft \downarrow \circlearrowright$$

n-morphisms



Yoga: reflective subcategories

Def $\mathcal{C} \text{ CAT} \rightarrow \hat{X}$ is S-local if $S \ni \begin{array}{ccc} A & \longrightarrow & C \\ \downarrow & \dashrightarrow & \uparrow \\ B & \cdots & ! \end{array}$

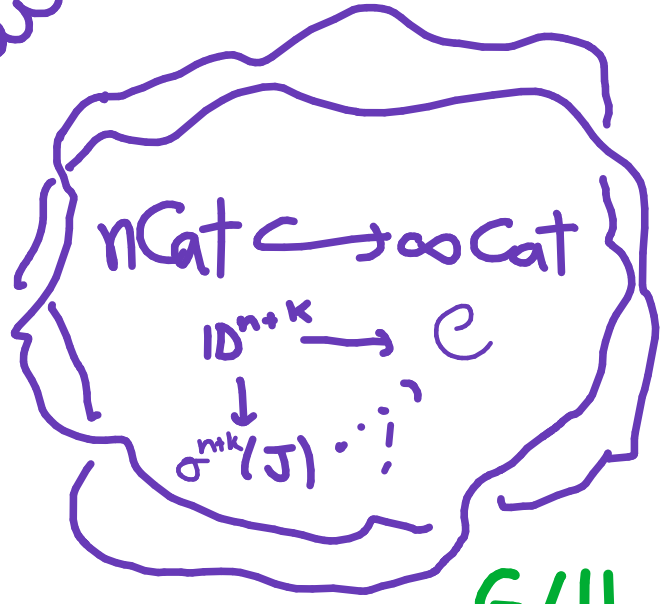
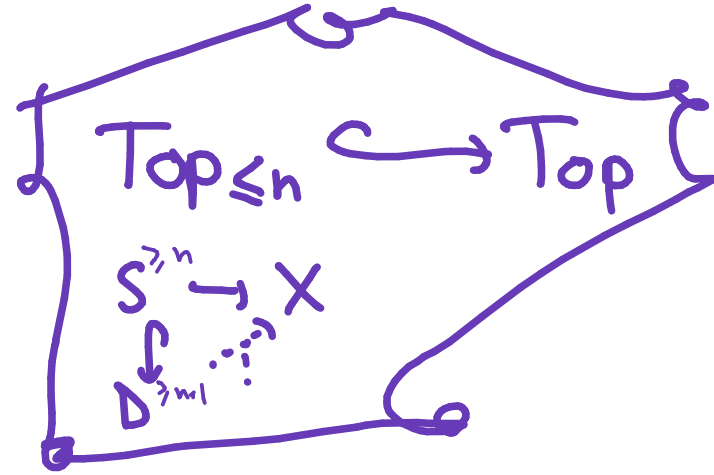
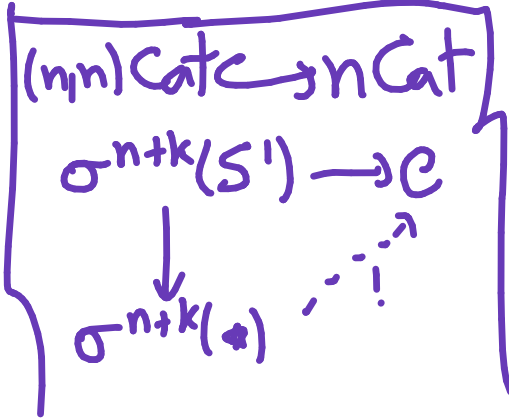
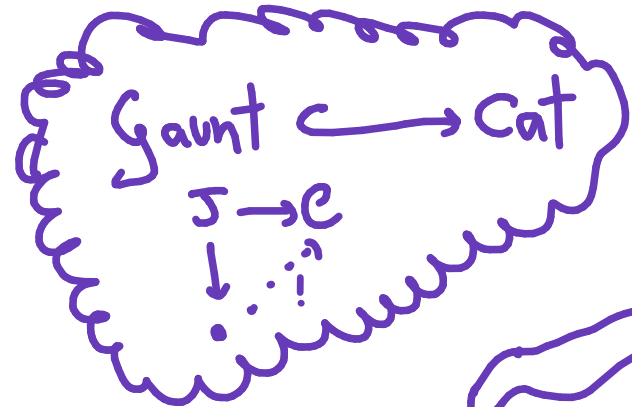
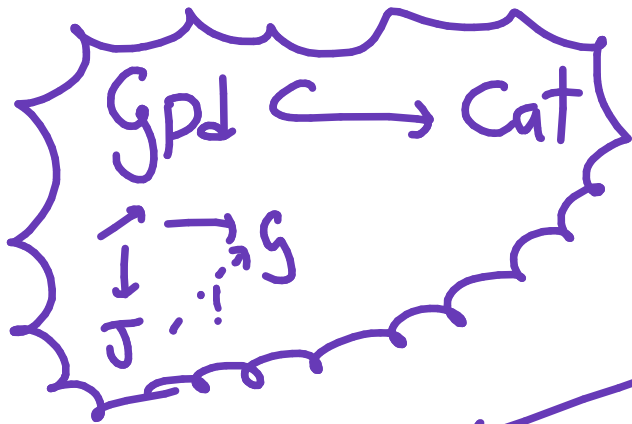
Prop \mathcal{C} presentable + $\text{dom}(S)$ compact

$\Rightarrow \mathcal{C}_S \xrightleftharpoons[L]{T} \mathcal{C}$ is a reflective subcategory

Sketch: $\hat{S} = \text{codiag. compl.}(S) \xrightarrow{\text{small obj arg}} (\text{coll}(\hat{S}), \text{urlp}(S)) \rightsquigarrow X \xrightarrow{\text{S-loc}} \hat{X} \xrightarrow{\text{urlp}(S)} \bullet$

Define $LX = \hat{X}$ 5/11

Yoga: reflective subcategories



Yoga: reflective subcategories

Def

$$\infty \text{Cat Adj} \begin{matrix} \xrightarrow{\tau} \\ \xleftarrow{\tau} \end{matrix} \infty \text{Cat}$$

$\sigma^n(\tau) \rightarrow e$
 \downarrow
 $\sigma^n(\text{adj}) \dots \overset{\tau}{!} \rightsquigarrow$ automatic bc epi

$\left\{ L(\mathbb{D}^n) \text{ is a conservative family} \right\}$

$\left\{ \mathbb{D}^n \right\}$ is a conservative family

or: $\infty \text{Cat Adj} = \varprojlim (\dots \rightarrow (n+1)\text{Cat Adj} \rightarrow n\text{Cat Adj} \rightarrow \dots)$

or: $\infty \text{Cat Adj} = \varprojlim (\dots \rightarrow \infty \text{Cat Adj}_n \rightarrow \infty \text{Cat Adj}_{n-1} \rightarrow \dots)$

Yoga: reflective subcategories

Def

$$\infty \text{Cat Adj} \begin{matrix} \xrightarrow{\tau} \\ \xleftarrow{L} \end{matrix} \infty \text{Cat}$$

"adjoining a lot of adjoints"

$\sigma^n(\tau) \rightarrow \mathcal{C}$
 \downarrow
 $\sigma^n(\text{adj}) \rightarrow \mathcal{C}$
 ! automatic bc epi

$\{L(\vec{D}^n)\}$ is a conservative family

$\neq \sigma(\text{adj})$

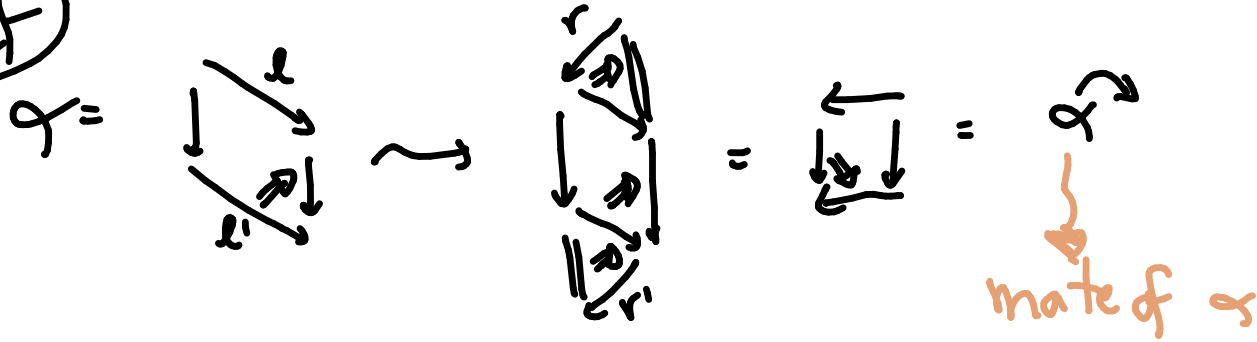
$\{\vec{D}^n\}$ is a conservative family

or: $\infty \text{Cat Adj} = \varprojlim (\dots \rightarrow (n+1) \text{Cat Adj} \rightarrow n \text{Cat Adj} \rightarrow \dots)$

or: $\infty \text{Cat Adj} = \varprojlim (\dots \rightarrow \infty \text{Cat Adj}_n \rightarrow \infty \text{Cat Adj}_{n-1} \rightarrow \dots)$

Some calculus, mate

Def



mate of α via Gray products

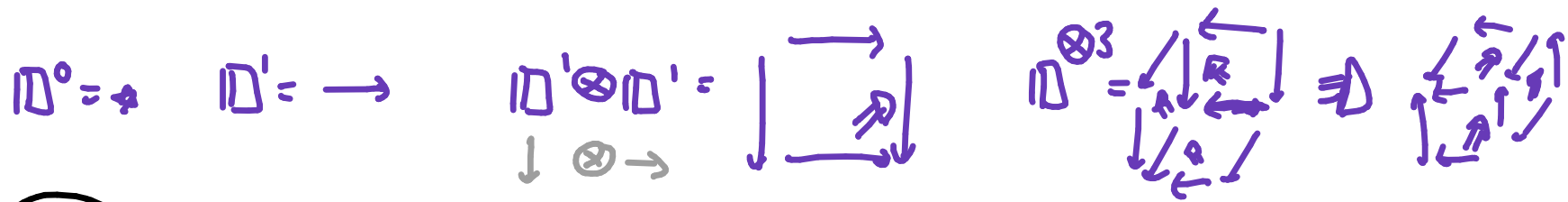
Def

α is Beck-Chevalley if

1. $\exists \alpha'$
2. $\exists (\alpha')^{-1}$

this w/ today (B.C. = $\mathbb{D}' \otimes \text{adj}$)

Some calculus, mate




Thm (Campion) $\exists!$ closed monoidal structure extending this:

$$- \otimes C : \infty\text{Cat} \rightleftarrows \infty\text{Cat} : \text{Fun}^{\text{lax}}(C, -) \quad \text{or} \quad C \otimes - \dashv \text{Fun}^{\text{oplax}}(C, D)$$

→ really hard to control outside of certain polygraph inclusions

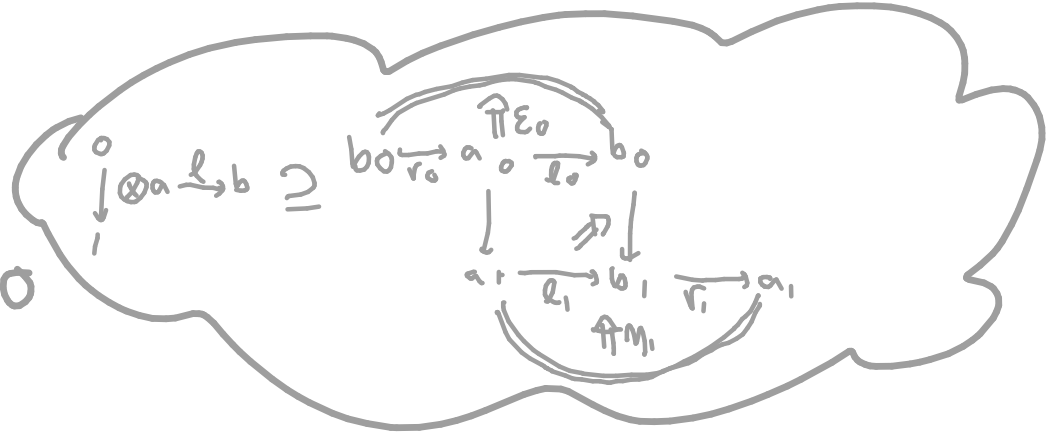
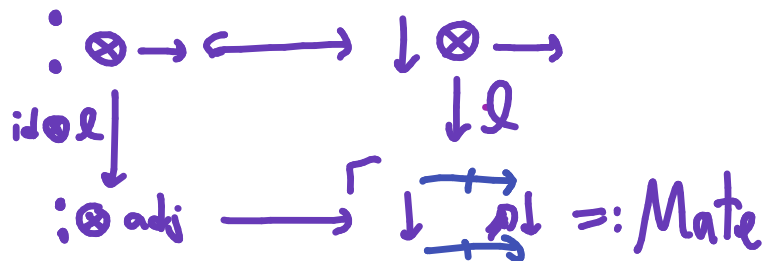
↳ just basic algebra here

e.g. $\mathbb{D}^2 \rightarrow \text{Fun}^{\text{lax}}(\mathbb{D}^1, C) \simeq \mathbb{D}^2 \otimes \mathbb{D}^1 \rightarrow C \simeq$


modifications

9/11

Walking mate



Def $\downarrow \otimes \rightarrow \xrightarrow{\text{lax n.t.}} \infty \text{ Cat}$
 $\downarrow \otimes \rightarrow \xrightarrow{\quad} \text{lax n.t.}$

Thm $\text{LFun}^{\text{lax}}(C, D) \simeq \text{RFun}^{\text{oplax}}(C, D)^{\text{coop}}$

Rmk $\mathbb{D}^1 \xrightarrow{\sim} \mathbb{D}^1{}^{\text{coop}}$
 $a \text{ adj} \xrightarrow{\sim} a \text{ adj}^{\text{coop}}$

$$\begin{array}{ccc}
 \downarrow \otimes \rightarrow & \xrightarrow{\sim} & \downarrow \otimes \rightarrow^{\text{coop}} \\
 \ell \downarrow & & \downarrow r \\
 \text{Mate} & \rightarrow & \text{Mate}^{\text{coop}}
 \end{array}$$

Walking mate

Thm $L\text{Fun}^{\text{lax}}(C, D) \simeq R\text{Fun}^{\text{oplax}}(C, D)^{\text{cop}}$

Conjecture (Marsden-Reutter) $\otimes |_{\infty\text{CatAdj}}$ is symmetric

Not a proof but.

Prop.(T.) If $C, D \in \infty\text{CatAdj}$, then $C \otimes D \simeq D \otimes C$.

Sketch: $L\text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{\text{lop, 2op}}$

fact: $\text{Fun}^{\text{lax}}(C, D) \in \text{Adj} \infty\text{Cat} \xleftrightarrow[\text{2-cells}]{\text{mate the}} \text{Fun}^{\text{oplax}}(C, D) = \text{Fun}^{\text{lax}}(C, D)^{\text{2op, 3op}}$

no time
sorry

Walking, mate

$$\boxed{\text{Thm}} \quad \text{LFun}^{\text{lax}}(C, D) \simeq \text{RFun}^{\text{oplax}}(C, D)^{\text{cop}}$$

Conjecture (Mazda-Reutter) $\otimes |_{\infty \text{CatAdj}}$ is symmetric

Not a proof but. $C \in \infty \text{Cat}$

Prop.(T.) IF $D \in \infty \text{CatAdj}$, then $C \otimes D \simeq D \otimes C$.

Sketch: $\text{LFun}^{\text{lax}}(C, D) = \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{\text{lop}, 2\text{op}}$

fact: $\text{Fun}^{\text{lax}}(C, D) \in \text{Adj} \infty \text{Cat} \xrightarrow[\text{2-cells}]{\text{mate th}} \text{Fun}^{\text{oplax}}(C, D) = \text{Fun}^{\text{lax}}(C, D)^{\text{2op}, 3\text{op}}$

(no time
sorry)

$\hookrightarrow \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{n, \text{nil op}} \forall n$, take $n \rightarrow \infty \Rightarrow \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)$
// //

$$-\otimes C \simeq C \otimes -$$

Walking, mate

Thm $L\text{Fun}^{\text{lax}}(C, D) \simeq R\text{Fun}^{\text{oplax}}(C, D)^{\text{cop}}$

→ "Cor": B, C bicategories w/ adjoints $\rightarrow B \otimes_{\text{lax}} C \simeq C \otimes_{\text{lax}} B$
 ↳ the technique works, but not the proof

Prop.(T.) If $\text{coadj DE} \infty \text{CatAdj}$, then $C \otimes D \simeq D \otimes C$.

Sketch: $L\text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{\text{lop, 2op}}$

fact: $\text{Fun}^{\text{lax}}(C, D) \in \text{Adj} \infty \text{Cat} \xrightarrow[\text{remate the 1-cells}]{\text{mate the 2-cells}} \text{Fun}^{\text{oplax}}(C, D) = \text{Fun}^{\text{lax}}(C, D)^{\text{2op, 3op}}$

↳ no time sorry

↳ $\text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{\text{n, nil op}} \forall n$, take $n \rightarrow \infty \Rightarrow \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)$

Walking, mate

Thm $L\text{Fun}^{\text{lax}}(C, D) \simeq R\text{Fun}^{\text{oplax}}(C, D)^{\text{cop}}$

→ Cor: B, C bicategories w/ adjoints $\rightarrow B \otimes_{\text{lax}} C \simeq C \otimes_{\text{lax}} B$

TY

Prop.(T.) If $C \in \text{Cat}$, $D \in \infty\text{CatAdj}$, then $C \otimes D \simeq D \otimes C$.

Sketch: $L\text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{\text{lop, 2op}}$

fact: $\text{Fun}^{\text{lax}}(C | \text{AdjCat}) \xrightarrow[\text{via mate, the 1-cells}]{\text{mate, the 2-cells}} \text{Fun}^{\text{oplax}}(C, D) = \text{Fun}^{\text{lax}}(C, D)^{\text{2op, 3op}}$

no time sorry

$\hookrightarrow \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)^{\text{n, nil op}} \forall n$, take $n \rightarrow \infty \Rightarrow \text{Fun}^{\text{lax}}(C, D) = \text{Fun}^{\text{oplax}}(C, D)$

$-\otimes C \simeq C \otimes -$