

# Some results for adjunctionable 2-categories

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March 9, 2026

It seems that the category of theory of  $(\infty, \infty)$ -categories with adjoints behave like the homotopy theory of spaces, except that we need to keep track of the 1-categories  $\Pi_n(\mathcal{C}) := h_1(\Omega^n \mathcal{C})$  instead of the sets  $\pi_n(X) := h_0(\Omega^n X)$ . In this note we give one example of such phenomena. There is no claim to originality: we are adapting the last theorem in [H25].

**Definition 1.** A morphism  $f : x \rightarrow y$  is a **left adjoint** if there exists a morphism  $g : y \rightarrow x$  and a pair of 2-morphisms

$$\begin{array}{c} & y & \\ f \nearrow & \Uparrow \eta & \searrow g \\ x & \xlongequal{\quad} & x \end{array} \quad \text{and} \quad \begin{array}{c} y & \xlongequal{\quad} & y \\ g \searrow & \Uparrow \varepsilon & \nearrow f \\ & x & \end{array} .$$

such that the *triangle identities* hold:

$$\begin{array}{c} & y & \xlongequal{\quad} & y \\ f \nearrow & \Uparrow \eta & \searrow g & \Uparrow \varepsilon & \nearrow f \\ x & \xlongequal{\quad} & x & & \end{array} = \begin{array}{c} & y & \xlongequal{\quad} & y \\ f \nearrow & 1_f & \nearrow f \\ x & \xlongequal{\quad} & x & \end{array}$$

and

$$\begin{array}{c} y & \xlongequal{\quad} & y \\ g \searrow & \Uparrow \varepsilon & \nearrow f & \Uparrow \eta & \searrow g \\ & x & \xlongequal{\quad} & x & \end{array} = \begin{array}{c} y & \xlongequal{\quad} & y \\ g \searrow & 1_g & \searrow g \\ & x & \xlongequal{\quad} & x & \end{array} .$$

**Example 2.** An object  $x$  in a monoidal category  $(\mathcal{M}, \otimes)$  is a left dual precisely if it is a left adjoint when regarded as a morphism in  $\mathcal{M}$ . △

**Proposition 3.** A morphism  $f : x \rightarrow y$  in  $\mathcal{C}$  is a left adjoint if and only if the induced functor  $f_* : \mathcal{C}(z, x) \rightarrow \mathcal{C}(z, y)$  is a left adjoint functor for each  $z \in \mathcal{C}$ , naturally in  $z$ .

*Sketch of proof.* If  $f \dashv f^R$ , then there is an induced adjunction  $f_* : \mathcal{C}(z, x) \rightleftarrows \mathcal{C}(z, y) : f_*^R$ , with unit and counit given by whiskering  $\eta$  and  $\varepsilon$ . Conversely, if  $f_* \dashv G$ , then the adjunction  $f \dashv f^R$  is obtained by setting  $f^R := G(1_y)$ . The unit is obtained by transposing  $1_f$  along  $f_* \dashv G$ :

$$1_f \in \mathcal{C}(x, y)(f, f) \simeq \eta \in \mathcal{C}(x, x)(1_x, f^R f)$$

The counit is obtained by transposing  $1_{f^R}$  to  $\varepsilon : f f^R \rightarrow 1_y$ , and the triangle identities follow from the triangle identities for  $f_* \dashv G$ . □

**Definition 4.** A 2-functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a **2-equivalence** if the induced functors

$$h_1 F : h_1 \mathcal{C} \rightarrow h_1 \mathcal{D} \quad \text{and} \quad F_{x,y} : \mathcal{C}(x, y) \rightarrow \mathcal{D}(Fx, Fy)$$

are equivalences of categories for every  $x, y \in \mathcal{C}$ .

The following result shows that the condition on  $F_{x,y}$  only needs to be checked on endomorphisms, provided that  $\mathcal{C}$  admits left adjoints.

**Theorem 5** (Heine). *Let  $\mathcal{C}$  be a 2-category that admits left adjoints. A 2-functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a 2-equivalence if and only if the induced functors*

$$h_1 F : h_1 \mathcal{C} \rightarrow h_1 \mathcal{D} \quad \text{and} \quad F_{x,x} : \mathcal{C}(x, x) \rightarrow \mathcal{D}(Fx, Fx)$$

are equivalences of categories for every  $x \in \mathcal{C}$ .

*Proof.* The forward direction is clear from Definition 4, so we show the converse.

Essential surjectivity for each  $F_{x,y}$  follows from the equivalence  $h_1 F$ , so it remains to show that  $F_{x,y}$  is fully faithful. In other words, that  $F_{x,y} : \mathcal{C}(x, y)(f, g) \rightarrow \mathcal{D}(Fx, Fy)(Ff, Fg)$  is a bijection for every  $f, g : x \rightarrow y$ .

Pick a right adjoint  $f \dashv f^R$ . The induced adjunction  $f_* : \mathcal{C}(x, y) \rightleftarrows \mathcal{C}(x, y) : f_*^R$  yields an equivalence

$$f_* : \mathcal{C}(x, y)(f, g) \simeq \mathcal{C}(x, x)(1_x, f^R f) : f_*^R.$$

Adjunctions are preserved by  $F$ , so  $Ff \dashv Ff^R$ , and hence the equivalence

$$Ff_* : \mathcal{D}(Fx, Fy)(Ff, Fg) \simeq \mathcal{D}(Fx, Fx)(1_{Fx}, Ff^R Ff) : Ff_*^R.$$

Under these identifications, the map  $F_{x,y}$  identifies with

$$\mathcal{C}(x, y)(f, g) \simeq \mathcal{C}(x, x)(1_x, f^R f) \xrightarrow{F_{x,x}} \mathcal{D}(Fx, Fx)(1_{Fx}, Ff^R Ff) \simeq \mathcal{D}(Fx, Fy)(Ff, Fg),$$

which is a bijection since  $F_{x,x}$  is an equivalence by hypothesis. □

**Corollary 6.** *In the situation of Theorem 5, if both categories are connected, then it suffices to check that  $F_{x,x}$  is an equivalence for a single  $x \in \mathcal{C}$ .*

## References

- [H25] Heine, H. (2025). *On the Categorification of Homology*. Preprint. Available at: <https://arxiv.org/abs/2505.22640v1>